EVENT CORRELATION SYSTEMS À LA CARTE

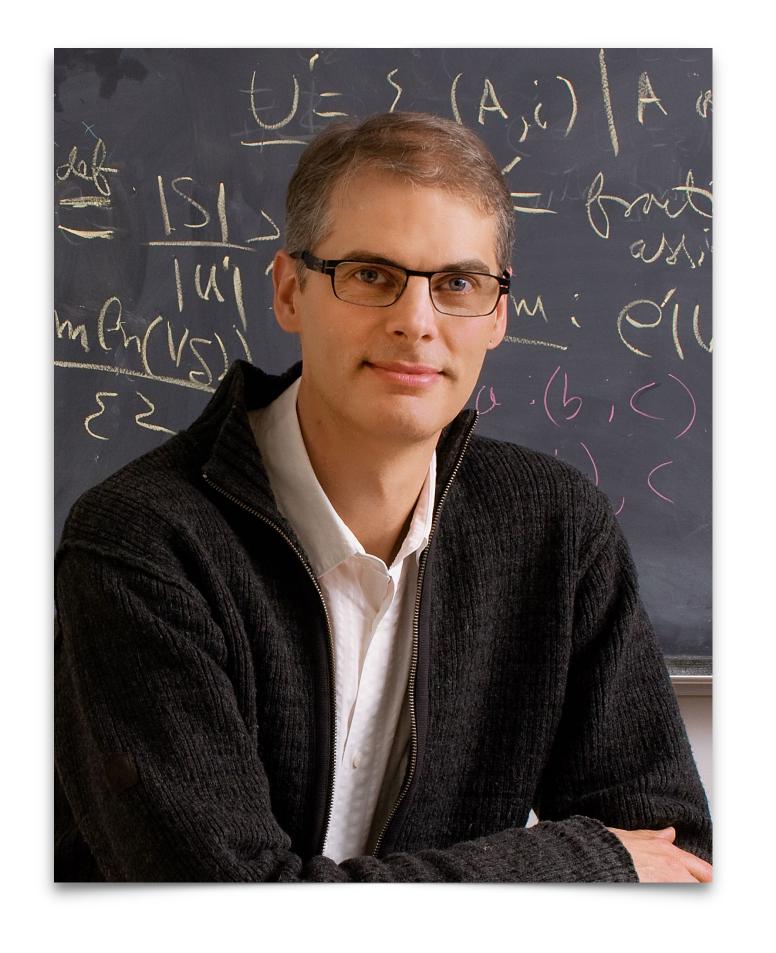
Oliver Bračevac

Software Technology Group @ TU Darmstadt, Germany

PurPL Seminar 2019-11-15

"PL people are good at taking complicated domains where people are creating new ways of doing things, but where it's too messy and not well understood, and pulling out the abstractions and the core concepts."

— Benjamin C. Pierce

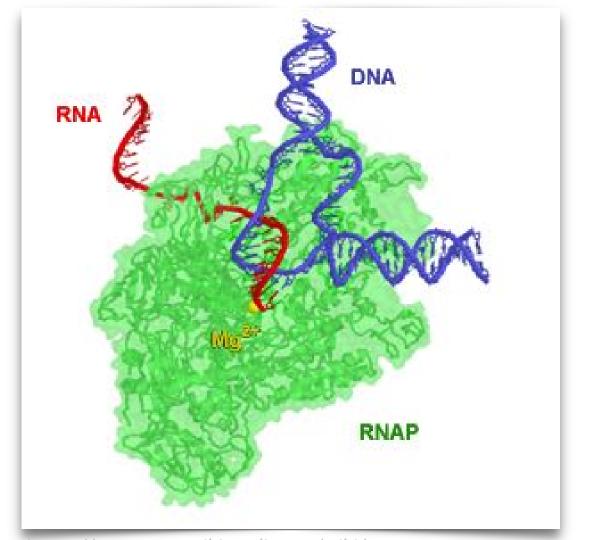


CONTEXT: EVENT CORRELATION

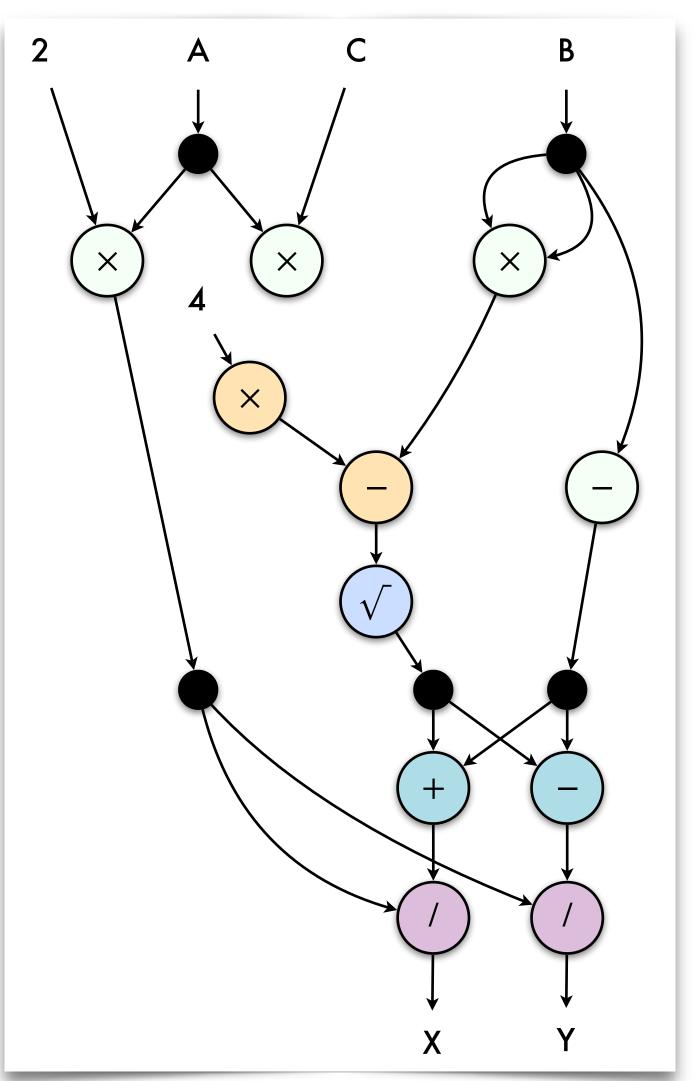
JOINING/MERGING OF INFORMATION FLOWS...

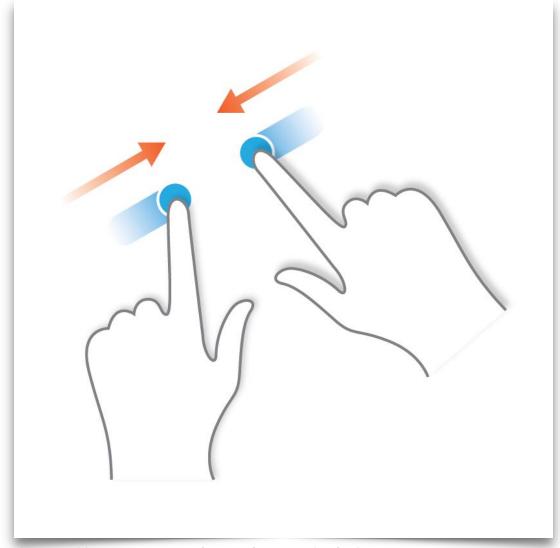


https://pxhere.com/en/photo/138038

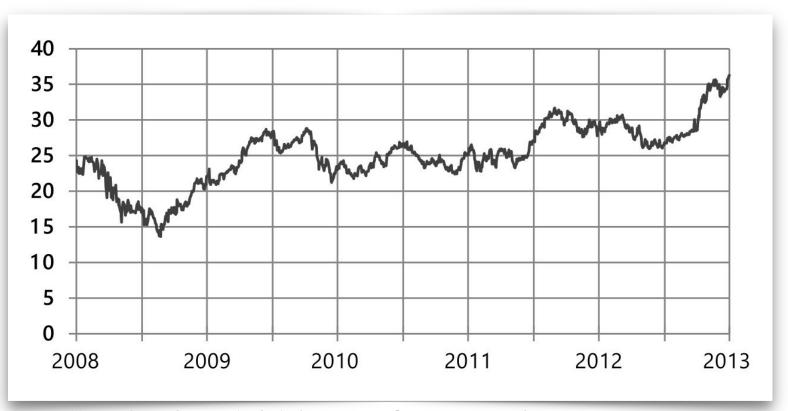


https://commons.wikimedia.org/wiki/File:RNAP_TEC_small.jpg





https://commons.wikimedia.org/wiki/ File:Gestures_Two_Hand_Pinch.png



https://en.wikipedia.org/wiki/File:Microsoft_5-Year_Stock_History.svg

CONTEXT: EVENT CORRELATION

... HAS BEEN REINVENTED SEVERAL TIMES OVER

Complex Event Processing

```
PATTERN SEQ(Stock+ a[], Stock b)
WHERE skip_till_next_match(a[], b) {
   [symbol]
   and a[1].volume > 1000
   and a[i].price > avg(a[..i-1].price)
   and b.volume < 0.8*a[a.LEN].volume }
WITHIN 1 hour
SASE+[Agrawal et al. '08])</pre>
```

Stream Processing

```
Select Istream(Close.item_id)
From Close[Now], Open[Range 5 Hours]
Where Close.item_id = Open.item_id
CQL[Arasu et al. '04]
```

Reactive Programming

```
val force = signal {
  mass() * accel()
}
ReScala [Salvaneschi et al. '14]
```

Concurrent Programming

```
def wait() & finished(r) =
  reply (Some r) to wait
or wait() & timeout() =
  reply None to wait

JoCaml [Conchon et al. '99]
```

CONTEXT: EVENT CORRELATION

... HAS BEEN REINVENTED SEVERAL TIMES OVER

```
Complex Event Processing
                                                   Reactive Programming
PATTERN SEQ(Stock+ a[], Stock b)
                                                   val force = signal {
                                                    mass() * accel()
WHERE skip_till_next_match(a[], b) {
  [symbol]
  and a[1].volume > 1000
                                                   ReScala [Salvaneschi et al. '14]
  and a[i].price > avg(a[..i-1].price)
  and b.volume < 0.8*a[a.LEN].volume
                        EVENT CORRELATION
WITHIN 1 hour
SASE+ [Agrawal et al. '08])
Stream Processing
                                                   Concurrent Programming
                                                   def wait() & finished(r) =
Select Istream(Close.item_id)
From Close[Now], Open[Range 5 Hours]
                                                     reply (Some r) to wait
Where Close.item_id = Open.item_id
                                                   or wait() & timeout() =
                                                     reply None to wait
CQL [Arasu et al. '04]
```

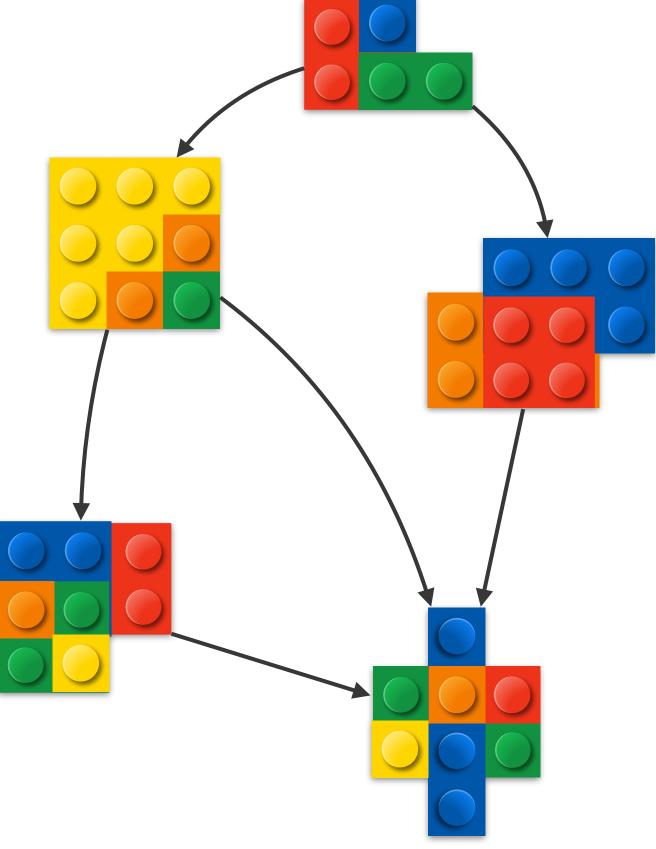
JoCaml [Conchon et al. '99]

MY THESIS IN A NUTSHELL

HOW DO WE OBTAIN EVENT CORRELATION SYSTEMS "À LA CARTE"?



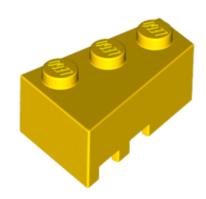
Freely composable correlation features



An extensible "feature menu"



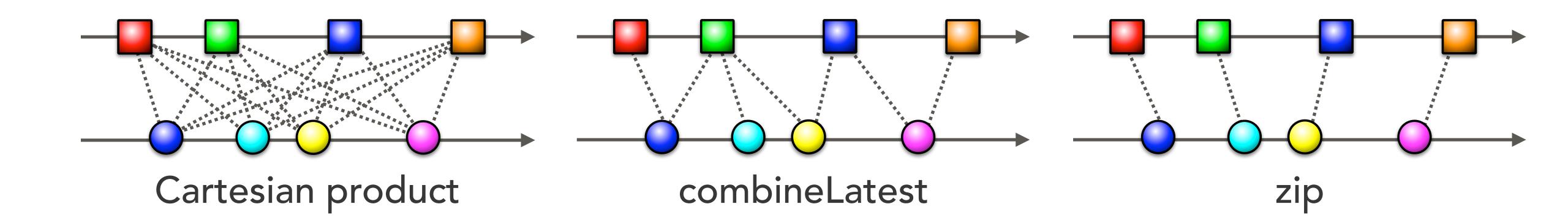




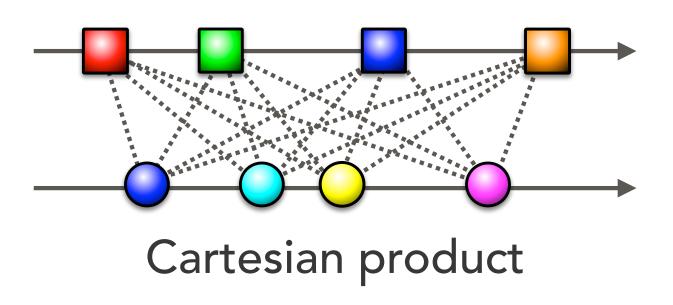
EXTENSIBLE CORRELATION SEMANTICS

THE ESSENCE OF JOINING

COMPUTATIONAL INTERPRETATION

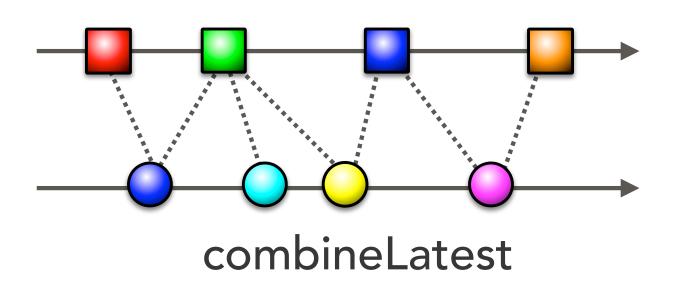


Enumerate the cartesian product, where (user-defined) side effects influence how the computation proceeds.



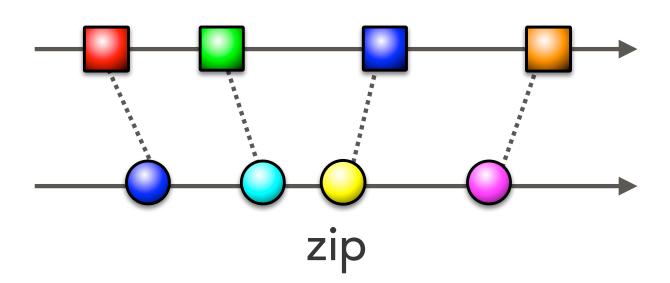
We support

Direct-style join specifications



We support

- Direct-style join specifications
- Controllable matching behavior



We support

- Direct-style join specifications
- Controllable matching behavior
- Mix-and-match composition of features
- Uniform programming model & extensibility

```
More in Thesis: Timing and Time Windows
correlate {
                                correlate {
                                  with slidingWindow(1h)
  x from r1
                                  x from r1
  y from r2
                                                              ns
  where
                                  y from r2
                                                             vior
    abs(?time_x - ?time_y) < 5s
                                  yield (x,y)}
  yield (x,y)}
                                       Uniform programming model &
  x from r1
                                      extensibility
  y from r2
 yield (x,y)}
```

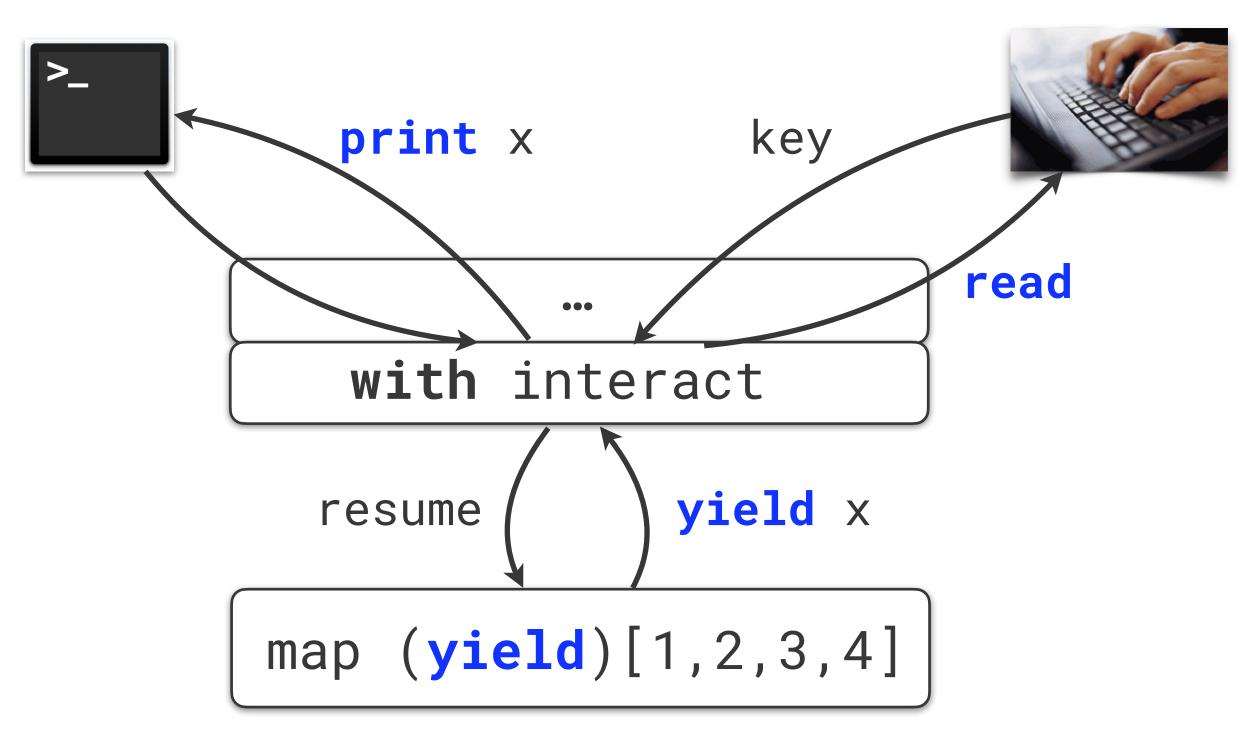
ALGEBRAIC EFFECTS AND HANDLERS

[PLOTKIN & POWER '03, PLOTKIN & PRETNAR '09]

Effect Interfaces:

yield: Int → Unit

Effect Handlers:



Effect Types:

interact: $\forall \alpha \mu. \{\alpha\}^{\langle \mathbf{yield}, \mu \rangle} \rightarrow^{\langle \mathbf{print}, \mathbf{read}, \mu \rangle}$ Unit

Handler composition:

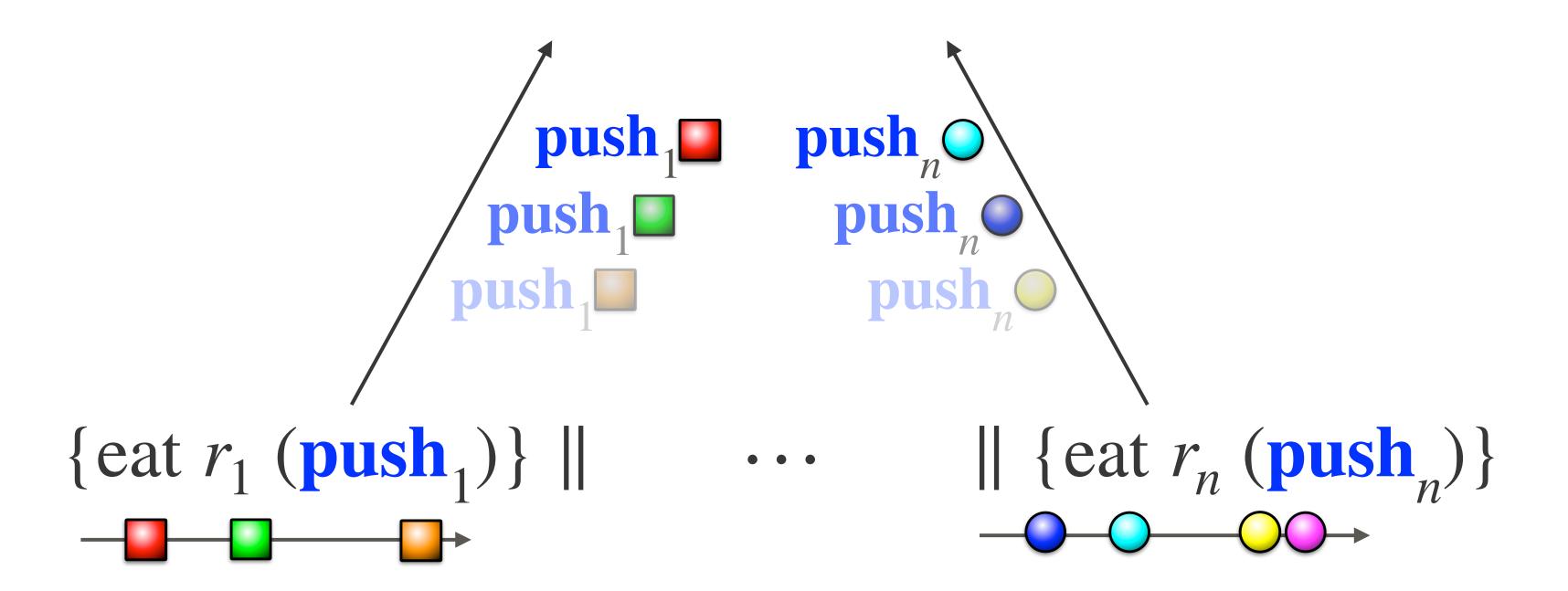
$$h_1 \coprod h_2 := \lambda tn \cdot \text{with } h_1 \text{ (with } h_2 \text{ (}tn \text{ (}\rangle\text{))}$$

DEFINES JOINS IN TERMS OF EFFECTS AND HANDLERS [BRACEVAC ET AL. 2018]

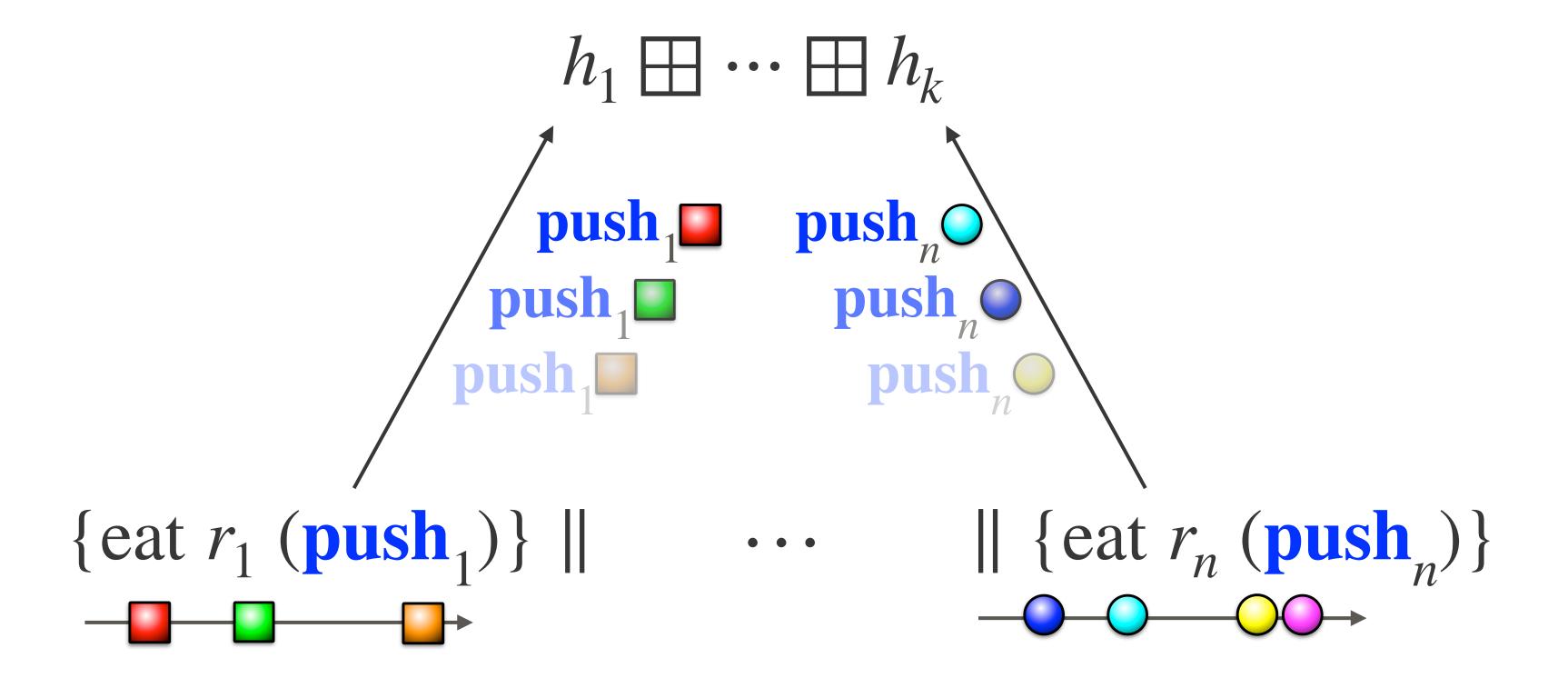
DEFINES JOINS IN TERMS OF EFFECTS AND HANDLERS [BRACEVAC ET AL. 2018]

$$\{\text{eat } r_1 \text{ (push}_1)\} \parallel \cdots \parallel \{\text{eat } r_n \text{ (push}_n)\}$$

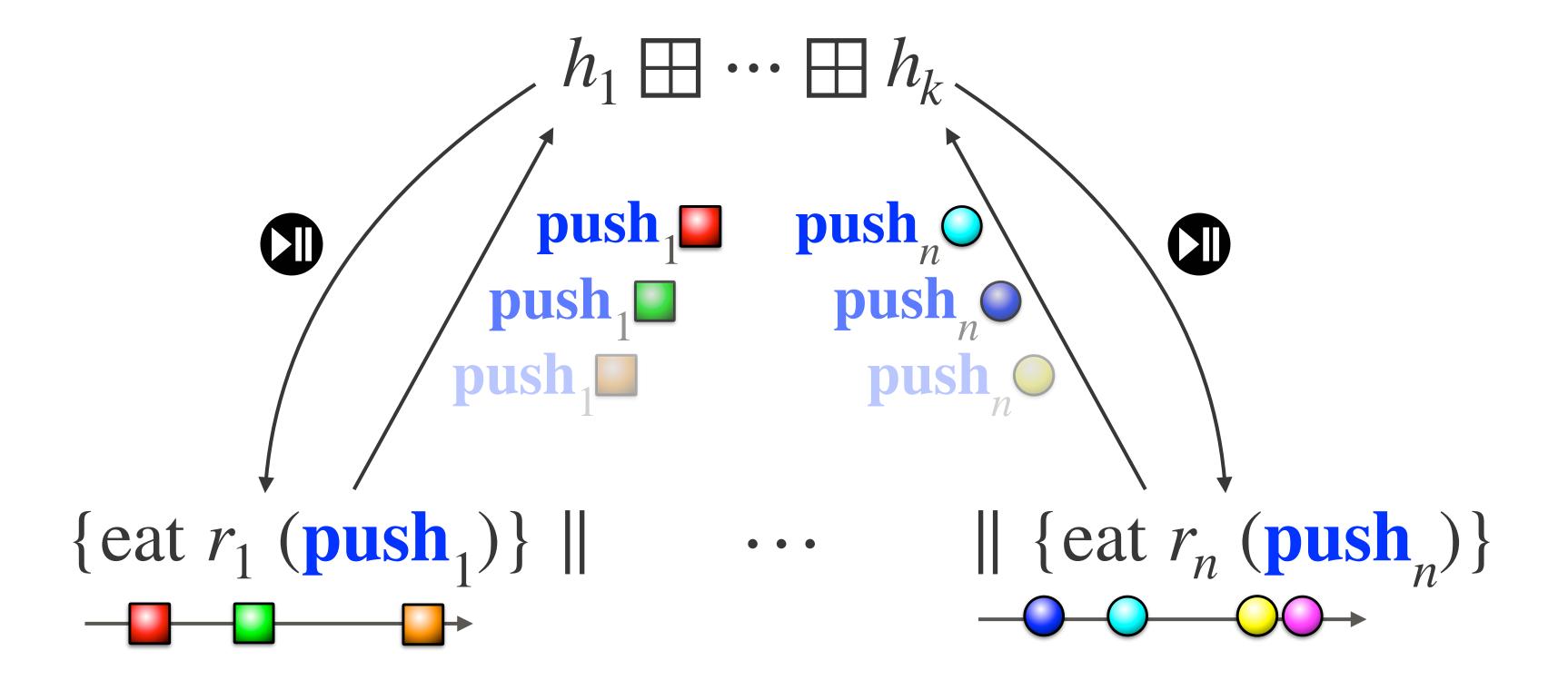
DEFINES JOINS IN TERMS OF EFFECTS AND HANDLERS [BRACEVAC ET AL. 2018]



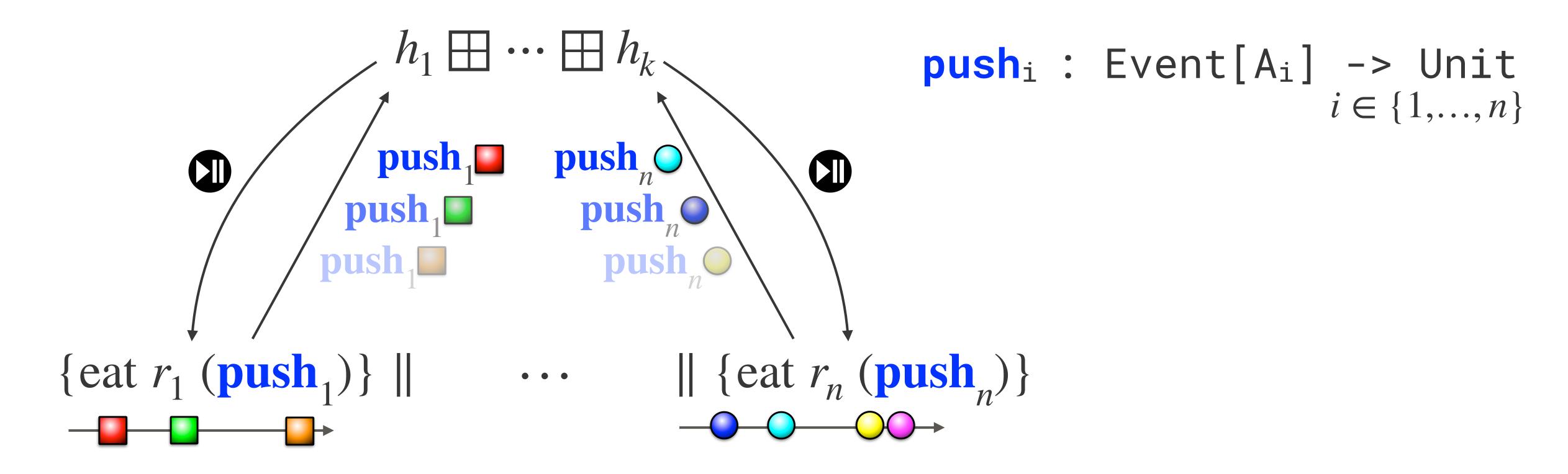
DEFINES JOINS IN TERMS OF EFFECTS AND HANDLERS [BRACEVAC ET AL. 2018]



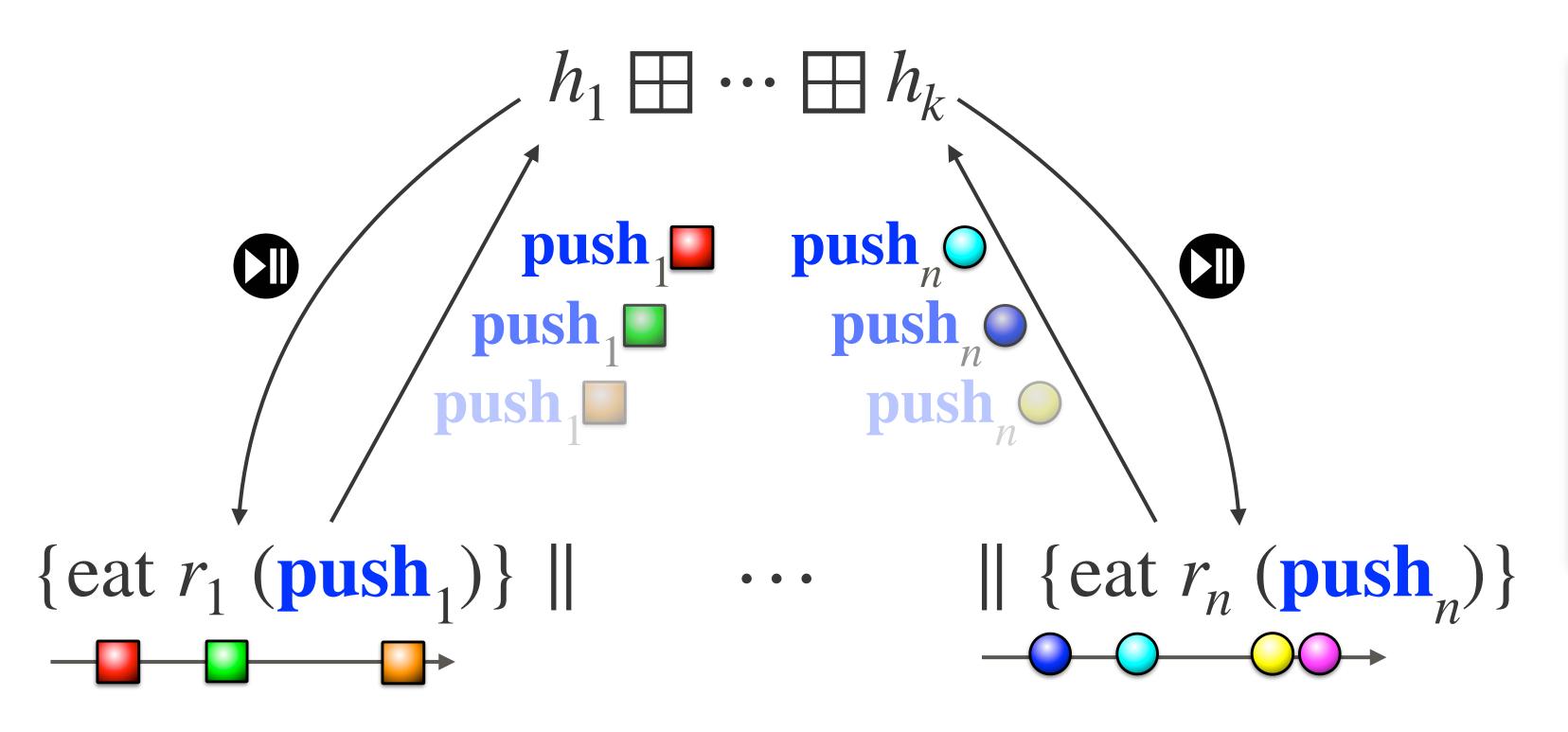
DEFINES JOINS IN TERMS OF EFFECTS AND HANDLERS [BRACEVAC ET AL. 2018]



DEFINES JOINS IN TERMS OF EFFECTS AND HANDLERS [BRACEVAC ET AL. 2018]

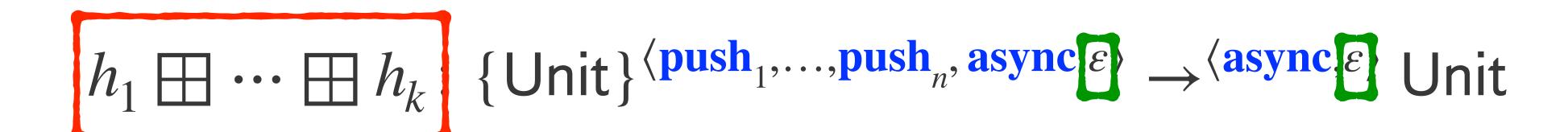


DEFINES JOINS IN TERMS OF EFFECTS AND HANDLERS [BRACEVAC ET AL. 2018]

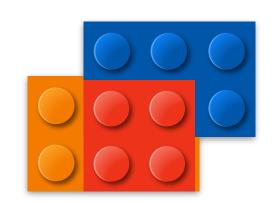


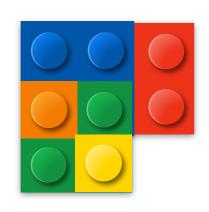


DEFINES JOINS IN TERMS OF EFFECTS AND HANDLERS [BRACEVAC ET AL. 2018]

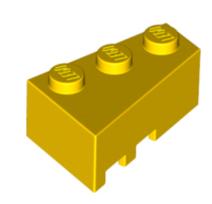


Legos!





New brick forms!



```
let rout =
  correlate {
    with h<sub>1</sub> ... ... h<sub>k</sub>
    x<sub>1</sub> from r<sub>1</sub>
    ...
    x<sub>n</sub> from r<sub>n</sub>
    where p
    yield e }
```

```
let r_{\text{out}} = \text{with } \left(h_{\otimes} \boxplus h_1 \boxplus \cdots \boxplus h_k\right)

\left\{\text{eat } r_1 \text{ (push}_1\right)\right\} \parallel \cdots \parallel \left\{\text{eat } r_n \text{ (push}_n\right)\right\}
```

```
let r_{\text{out}} =
correlate {

with h_1 \boxplus \cdots \boxplus h_k

x_1 from r_1

...

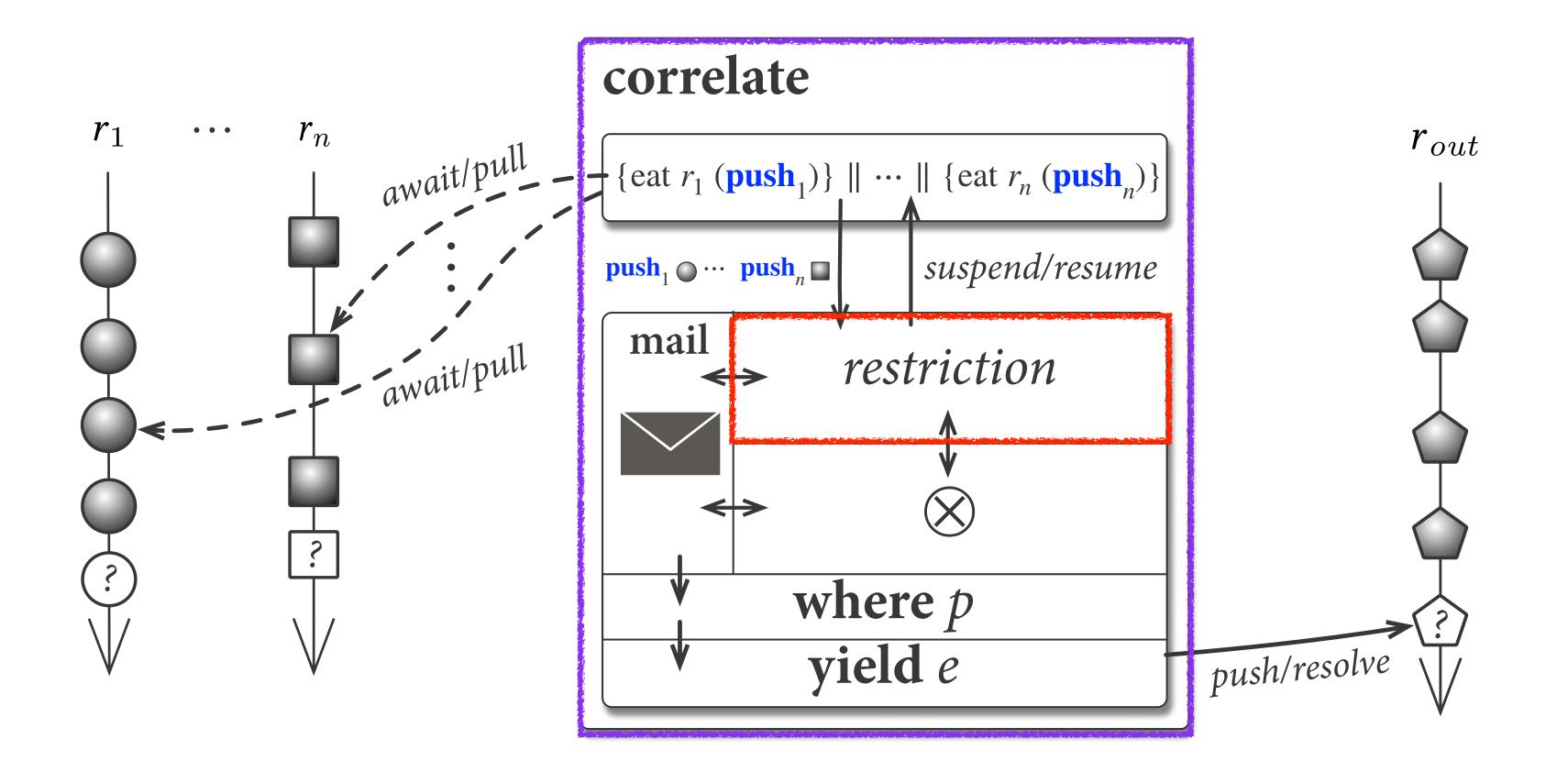
x_n from r_n

where p

yield e }

let r_{\text{out}} = \text{with } (h_{\otimes} \boxplus h_1 \boxplus \cdots \boxplus h_k)

{eat r_1 (push<sub>1</sub>)} \parallel \cdots \parallel {eat r_n (push<sub>n</sub>)}
```



```
let r_{\text{out}} =
correlate {

with h_1 \boxplus \dots \boxplus h_k

x_1 from r_1

\dots

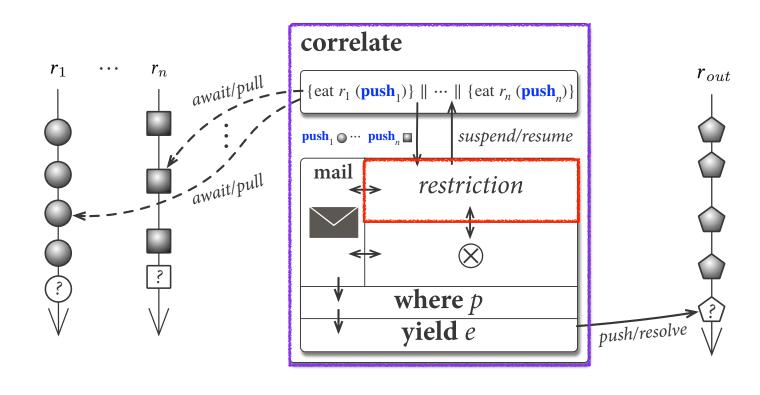
x_n from r_n

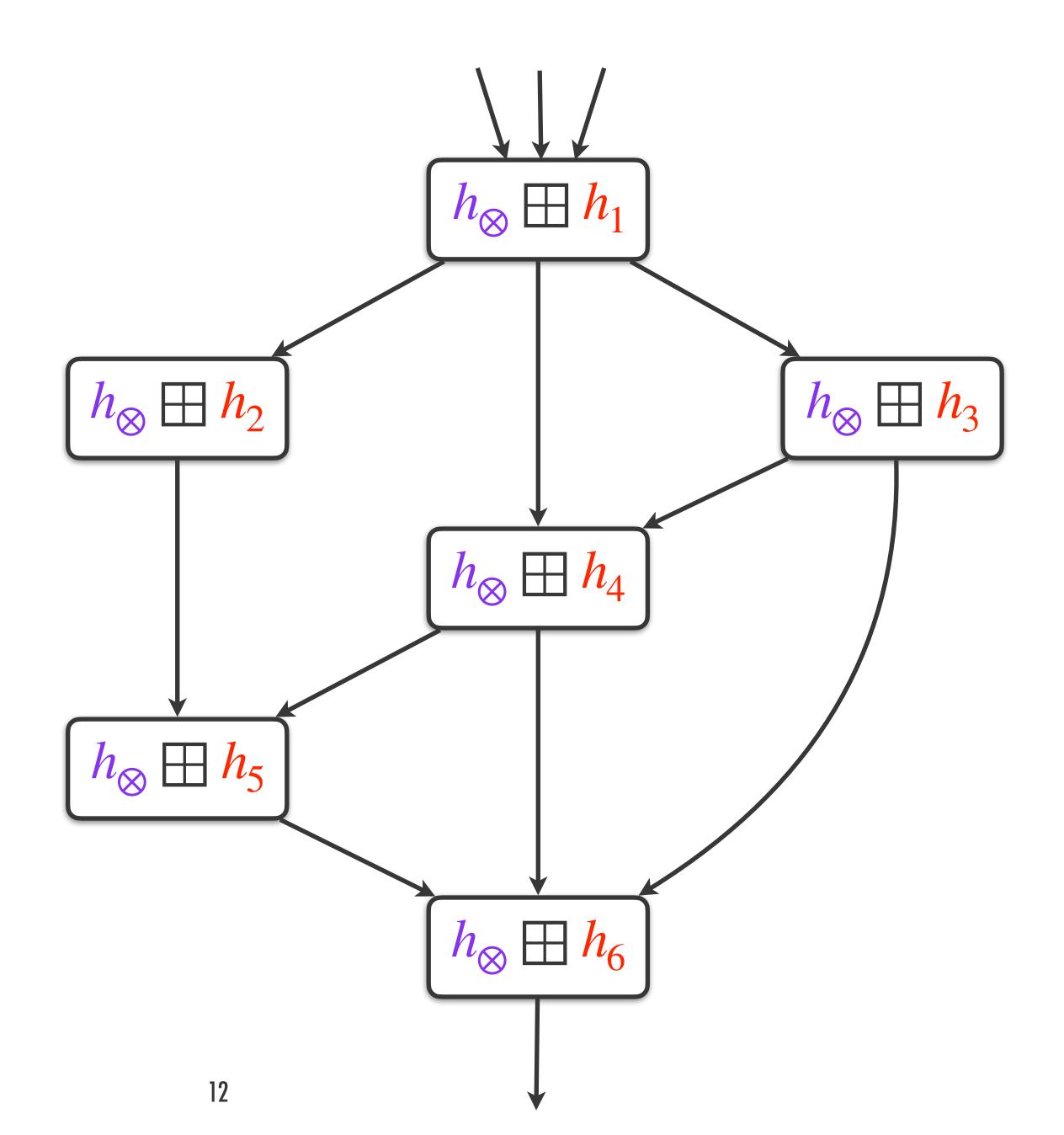
where p

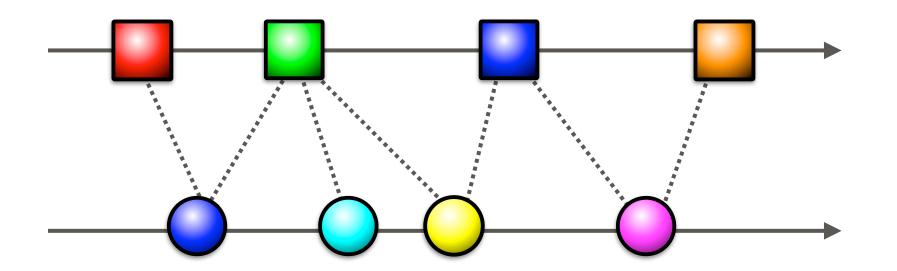
yield e }

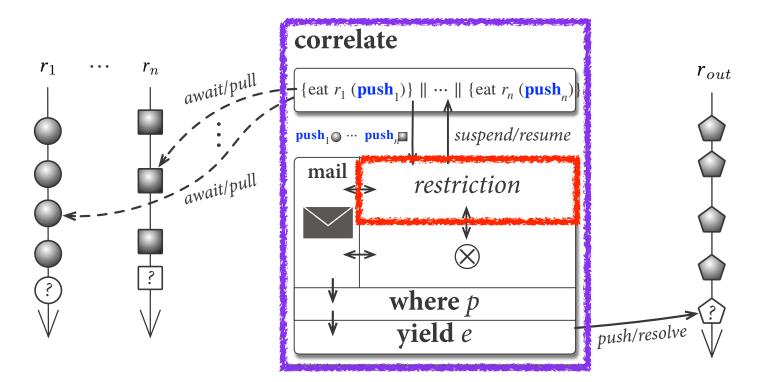
let r_{\text{out}} = \text{with } (h_{\otimes} \boxplus h_1 \boxplus \dots \boxplus h_k)

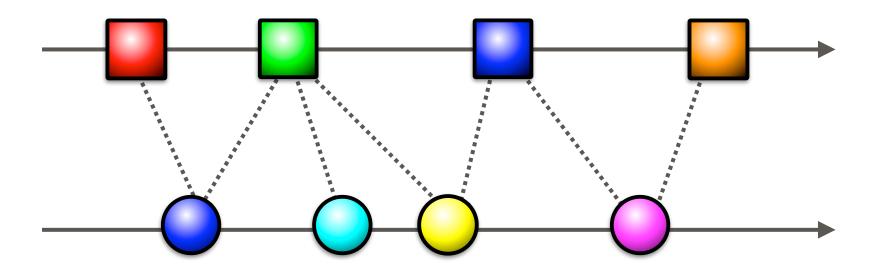
{eat r_1 (push<sub>1</sub>)} \parallel \dots \parallel {eat r_n (push<sub>n</sub>)}
```

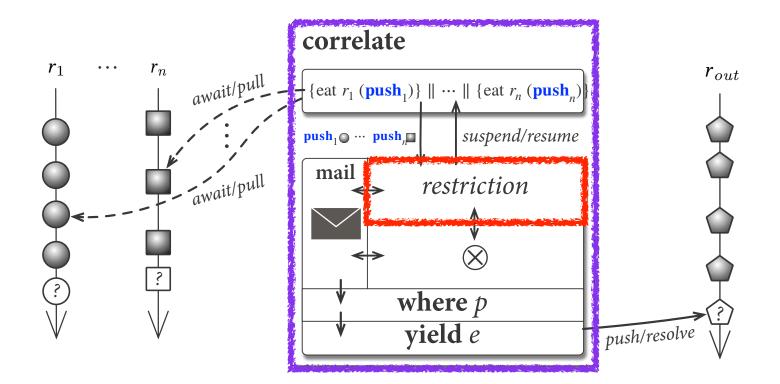




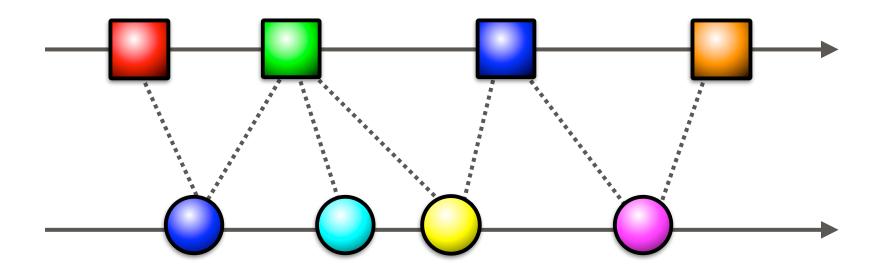


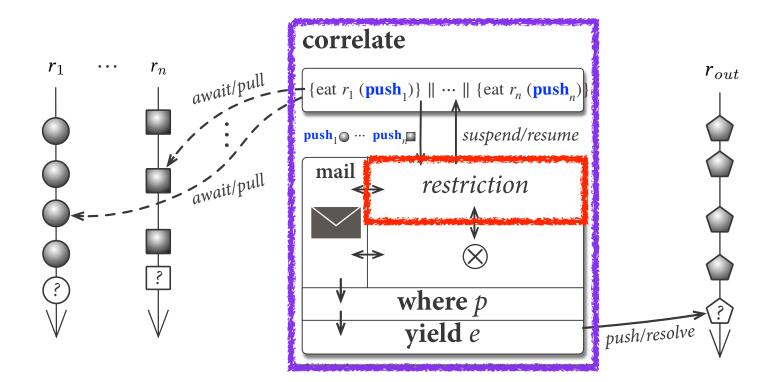




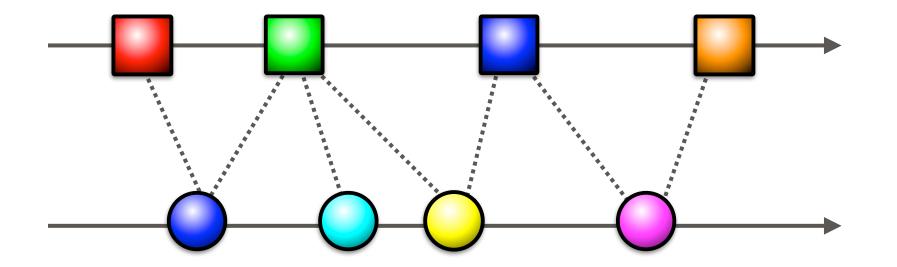


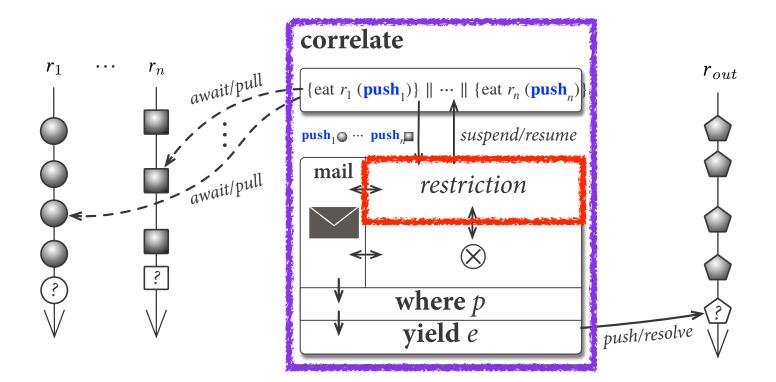
```
let mostRecentlyi = handler
| pushi ev resume ->
    seti [ev];
    resume (pushi ev)
```





```
let mostRecentlyi = handler
| pushi ev resume ->
seti [ev];
resume (pushi ev)
```





```
let mostRecentlyi = handler
| pushi ev resume ->
    seti [ev];

resume (pushi ev)
```

IMPLEMENTATION CHALLENGE: HOW CAN WE CORRELATE ANY NUMBER OF EVENT SOURCES?

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Declarative Pattern syntax

```
correlate {
   with h
   x<sub>1</sub> from r<sub>1</sub> ... x<sub>n</sub> from r<sub>n</sub>
   where p
   yield e }
```

IMPLEMENTATION CHALLENGE: HOW CAN WE CORRELATE ANY NUMBER OF EVENT SOURCES?

```
Declarative Pattern syntax
```

```
correlate { with h x_1 from x_1 ... x_n from x_n where p x_n \in \mathbb{N} yield e }
```

IMPLEMENTATION CHALLENGE: HOW CAN WE CORRELATE ANY NUMBER OF EVENT SOURCES?

```
Declarative Pattern syntax
```

```
correlate {
   with h
   x<sub>1</sub> from r_1 ... x_n from r_n
   where p
   yield e }
```

Vs. Comprehensions? LINQ [Meijer et al. '06]?

```
from x_1 in r_1 ... from x_n in r_n where p select e
```

IMPLEMENTATION CHALLENGE: HOW CAN WE CORRELATE ANY NUMBER OF EVENT SOURCES?

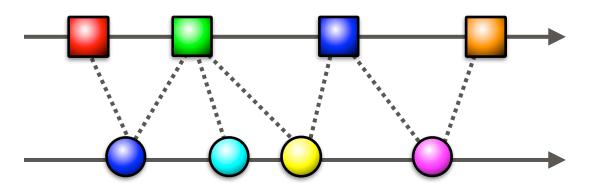
Declarative Pattern syntax

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correlate { with h x_1 from x_1 ... x_n from x_n where p x_n \in \mathbb{N} yield e }
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from x_1 in r_1 ... from x_n in r_n where p select e
```

Cannot express dataflow!



IMPLEMENTATION CHALLENGE: HOW CAN WE CORRELATE ANY NUMBER OF EVENT SOURCES?

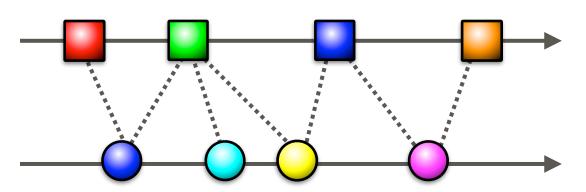
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Generative effect declarations...

```
effect push<sub>i</sub>: Event[A<sub>i</sub>] -> Unit i \in \{1,...,n\}
```

IMPLEMENTATION CHALLENGE: HOW CAN WE CORRELATE ANY NUMBER OF EVENT SOURCES?

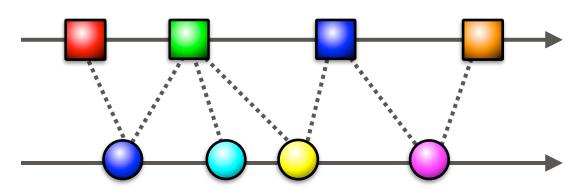
```
Declarative Pattern syntax
```

```
correlate { with h x_1 from x_1 ... x_n from x_n where p x_n \in \mathbb{N} yield e }
```

Vs. Comprehensions? LINQ [Meijer et al. '06]?

```
from x_1 in r_1 ... from x_n in r_n where p select e
```

Cannot express dataflow!



Generative effect declarations...

```
effect push<sub>i</sub> : Event[A<sub>i</sub>] -> Unit i \in \{1,...,n\}
```

...and their handlers...

```
let mostRecently<sub>i</sub> = handler
| push<sub>i</sub> ev resume ->
    set<sub>i</sub> [ev];
    resume (push<sub>i</sub> ev)
i \in \{1, ..., n\}
```

FROM THEORY TO PRACTICE

IMPLEMENTATION CHALLENGE: HOW CAN WE CORRELATE ANY NUMBER OF EVENT SOURCES?

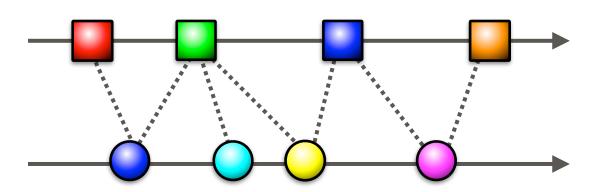
```
Declarative Pattern syntax
```

```
correlate { with h x_1 from x_1 ... x_n from x_n where p x_n \in \mathbb{N} yield e }
```

Vs. Comprehensions? LINQ [Meijer et al. '06]?

```
from x_1 in r_1 ... from x_n in r_n where p select e
```

Cannot express dataflow!



Generative effect declarations...

```
effect push<sub>i</sub> : Event[A<sub>i</sub>] -> Unit i \in \{1, ..., n\}
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...and their handlers...

```
let mostRecently<sub>i</sub> = handler
  | push<sub>i</sub> ev resume ->
    set<sub>i</sub> [ev];
    resume (push<sub>i</sub> ev)
    i \in \{1,...,n\}
```

Informal, lack linguistic support!

EXTENSIBLE PATTERN SYNTAX

A BETTER LINQ, WITHOUT COMPILER EXTENSIONS [BRACEVAC ET AL. 2020]

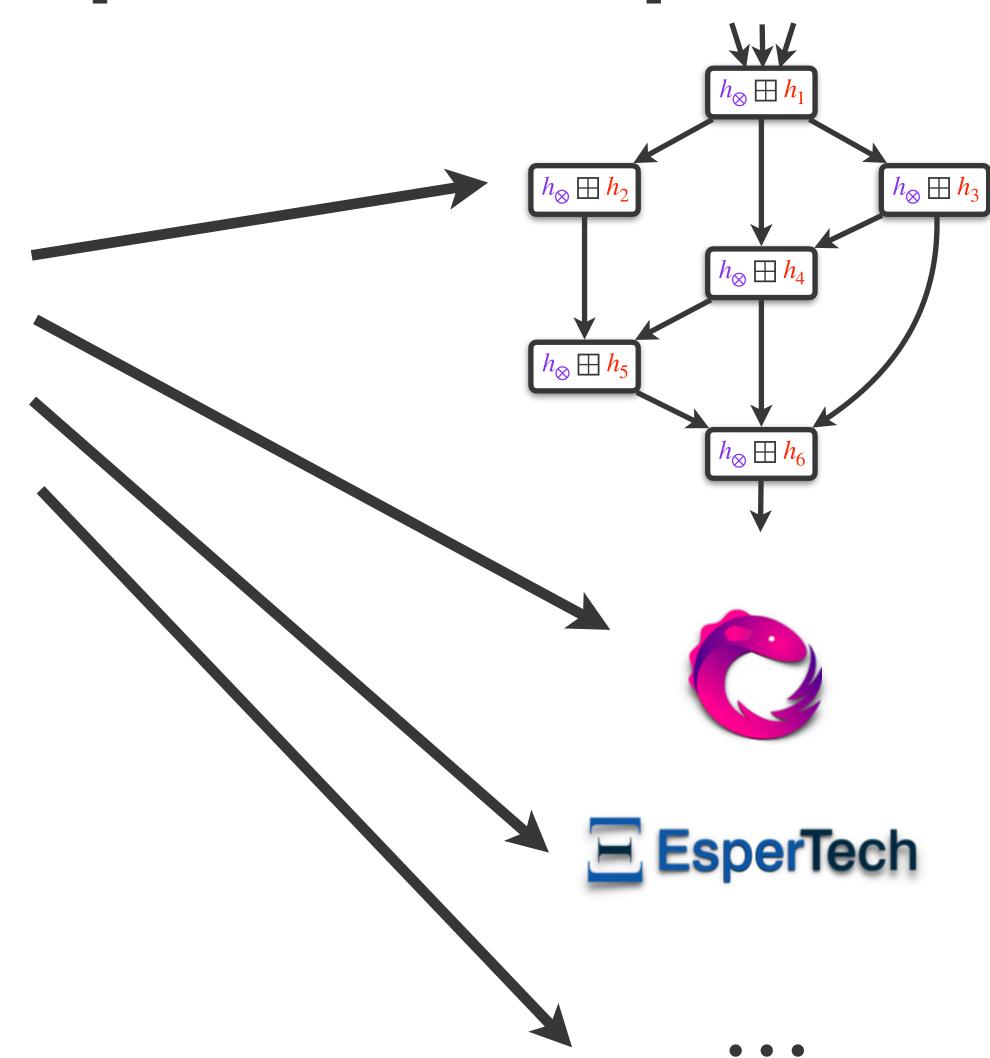
Portable!











POLYVARIADICITY

ARITY GENERICITY + HETEROGENEITY

Pattern Variables in Higher-Order Abstract Syntax (HOAS)

```
correlate: ('a, 'b) ctx -> ('b -> 'c pat) -> 'c shape repr
```

End-user query determines the type-level instantiations, i.e.,

```
correlate ((from temp_sensor) @. (from smoke_sensor))
```

Compiler calculates number and types of Variables automatically:

```
((float repr * meta repr) * (bool repr * meta repr) -> 'c pat) -> 'c shape repr
```

POLYVARIADICITY

ARITY GENERICITY + HETEROGENEITY

Pattern Variables in Higher-Order Abstract Syntax (HOAS)

```
correlate: ('a, 'b) ctx -> ('b -> 'c pat) -> 'c shape repr
```

Indexing by type vectors (contexts): $\mathbf{a} \equiv \overrightarrow{A} \in \mathsf{Type}^*$

End-user query determines the type-level instantiations, i.e.,

```
correlate ((from temp_sensor) @. (from smoke_sensor))
```

Compiler calculates number and types of Variables automatically:

```
((float repr * meta repr) * (bool repr * meta repr) -> 'c pat) -> 'c shape repr
```

```
correlate: ('a, 'b) ctx -> ('b -> 'c pat) -> 'c shape repr
```

```
correlate: ('a, 'b) ctx -> ('b -> 'c pat) -> 'c shape repr
```

```
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| pushi ev resume ->
    seti [ev];
    resume (pushi ev)
```

```
correlate: ('a, 'b) ctx -> ('b -> 'c pat) -> 'c shape repr
```

```
correlate: ('a, 'b) ctx -> ('b -> 'c pat) -> 'c shape repr
Context Polymorphism
let mostRecently: type a i. (i, a) ptr -> a ext
  fun ptr ctx ->
  let module I = project (ptr ()) ctx in
  handler
     I.push ev resume ->
     I.set [ev];
     resume (I.push ev)
```

```
correlate: ('a, 'b) ctx -> ('b -> 'c pat) -> 'c shape repr
```

```
correlate: ('a, 'b) ctx -> ('b -> 'c pat) -> 'c shape repr
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correlate: ('a, 'b) ctx -> ('b -> 'c pat) -> 'c shape repr
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correlate: ('a, 'b) ctx -> ('b -> 'c pat) -> 'c shape repr
let mostRecently: type a i. (i, a) ptr -> a ext
 fun ptr ctx ->
 let module I = project (ptr ()) ctx in
 handler
  | I.push ev resume ->
   I.set [ev];
   resume (I.push ev)
correlate ((from temp_sensor) @. (from smoke_sensor))
               (mostRecently 0) \boxplus (mostRecently 8)
               (fun ((temp, t1), ((smoke, t2), ())) ->
                    where p
                     (yield e))
```

III EVALUATION

HOW "EFFECTIVE" ARE RESTRICTION HANDLERS?

ZIP

60000000

• Experiment: Uniformly distribute 10⁷ events over 3 push streams.

MOSTRECENT

611

3-WAY CART.

525

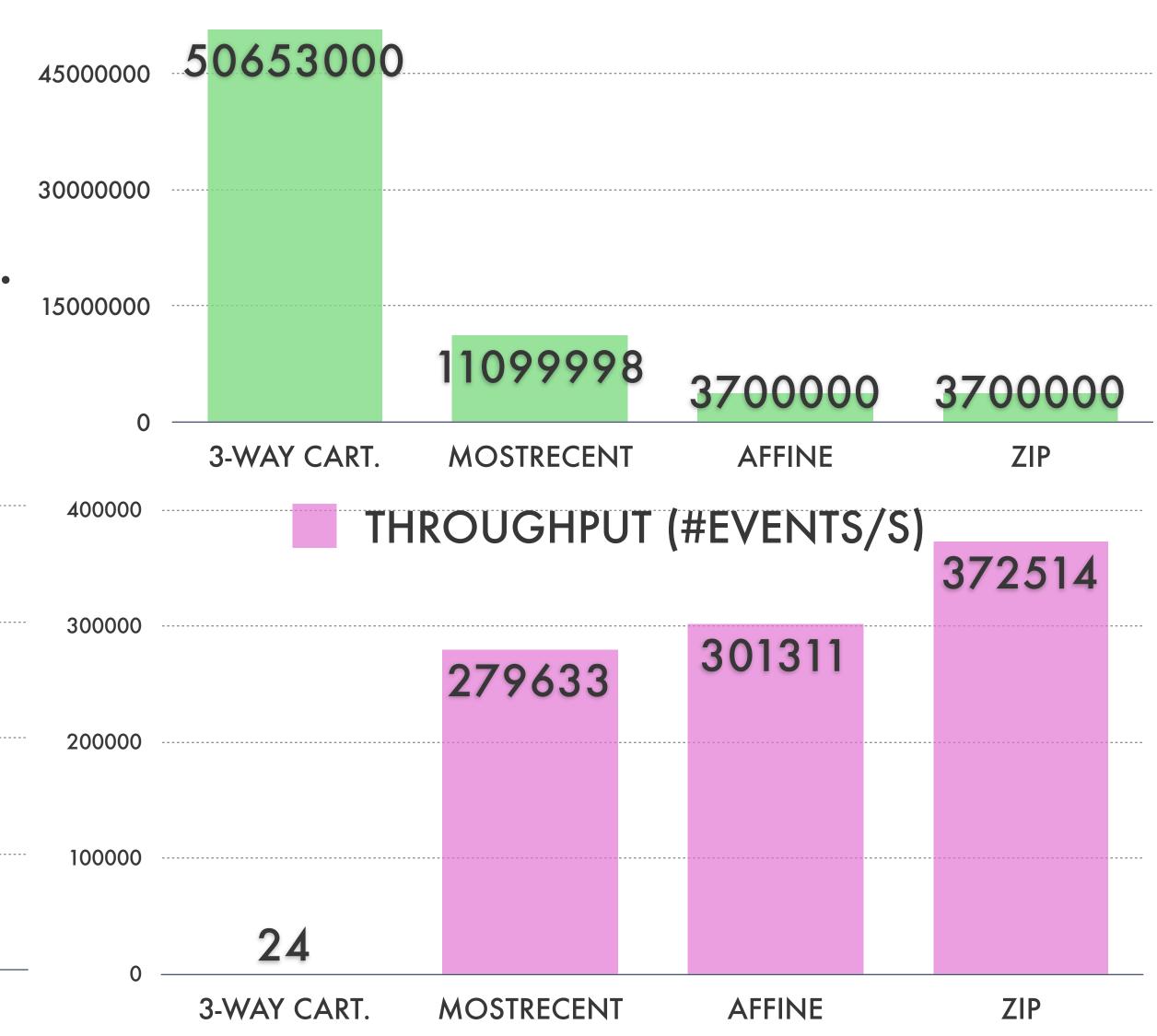
350

175

• Decomposition into naive cartesian product and restrictions is practical.

MAILBOX (#EVENTS/100 PUSHES)

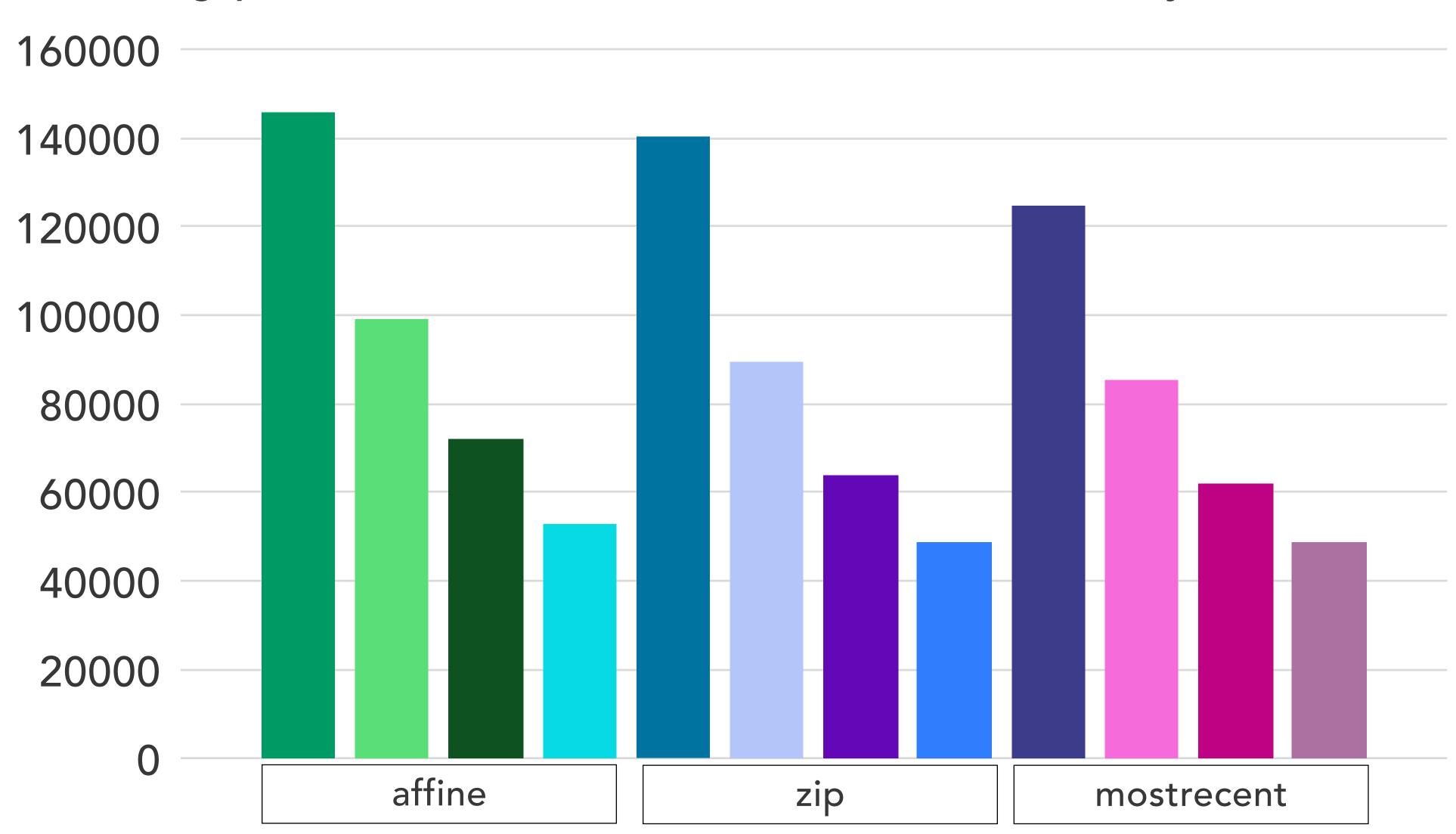
AFFINE



#MATERIALIZED TUPLES

EFFECT OF ARITY ON PERFORMANCE

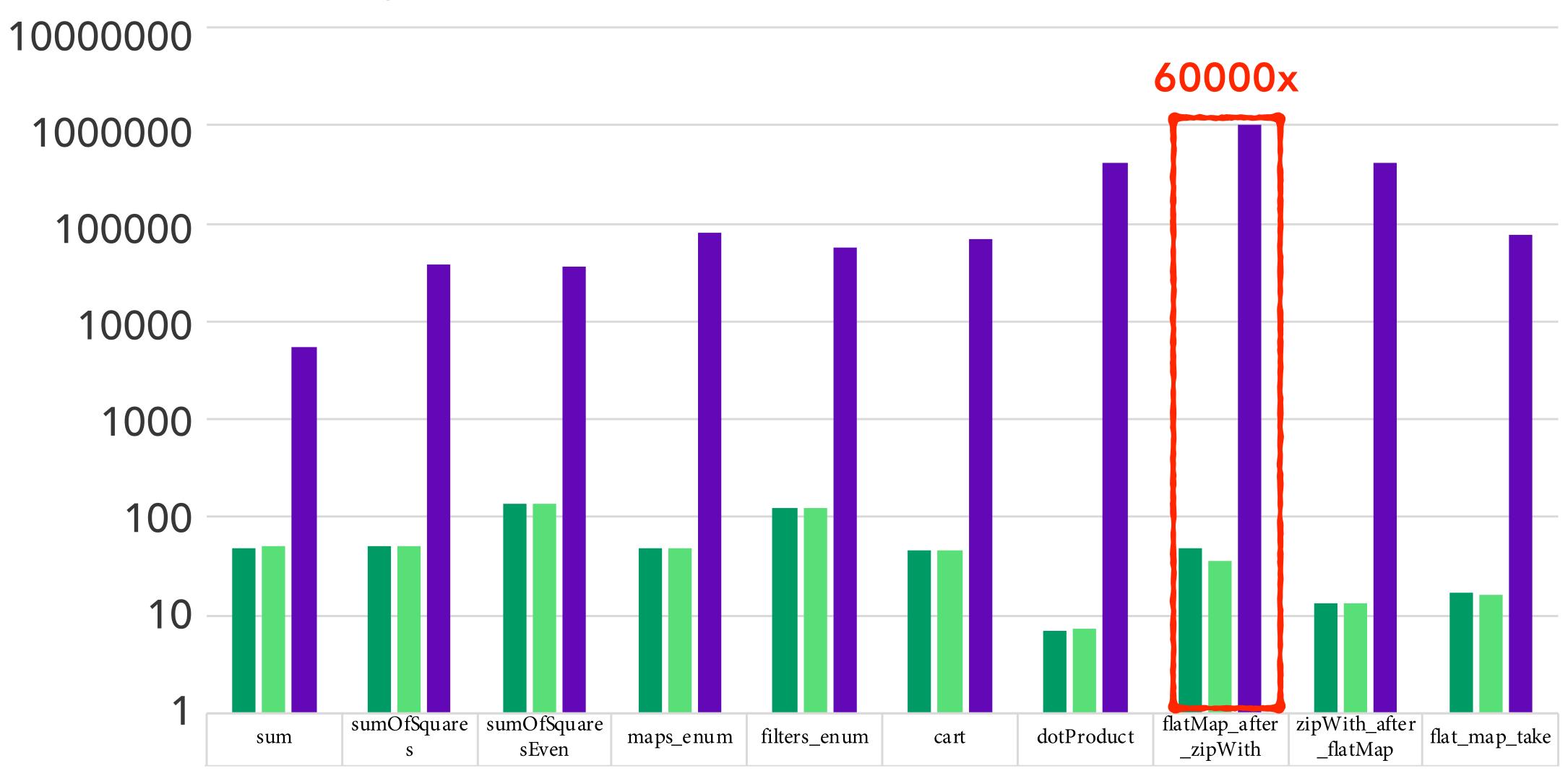
Total Throughput (#Events/sec), Even Distribution of 10^7 Events, Arity n = 8,12,16,20



COMPARISON TO HAND-WRITTEN CODE

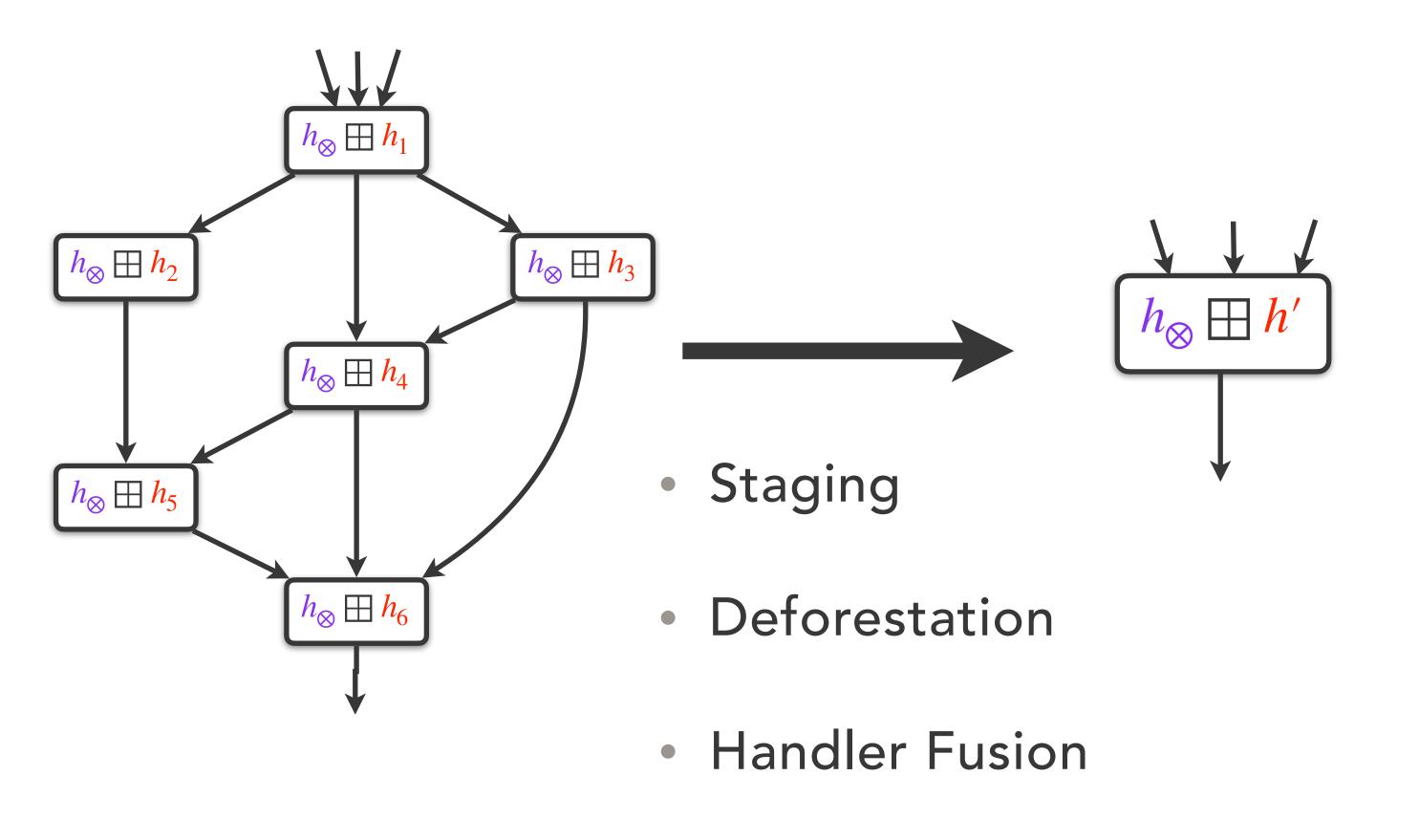
STRYMONAS LIBRARY [KISELYOV ET AL. 2017]

Execution Time (ms), log scale



FUTURE WORK

Make it fast!



Logical Interpretation
Propositions as Types, Proofs as Programs

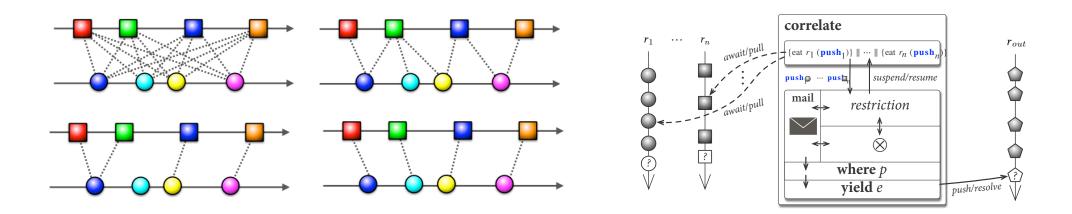
$$R[T] = \square (T \times \lozenge[R[T]])$$

- Strong Normalization
- (Co)Induction
- Liveness & Productivity

MY CONTRIBUTIONS

AN À LA CARTE SYNTAX AND SEMANTICS FOR EVENT CORRELATION

CARTESIUS: SEMANTIC MODEL



- The **first** work that **uniformly** expresses join variants.
- Programming of "interleaved coroutines".
- Effects make naïve enumerations efficient.

POLYJOIN: PORTABLE PL SYNTAX EMBEDDING

```
correlate ((from s_1) @. ... @. (from s_n)) (fun (x_1, (... (x_n, ())...) -> (where p (yield e)))
```

- A better LINQ: Practical, portable, extensible, type-safe correlation patterns.
- Join syntax and semantics are polyvariadic.

PROCESSING MODEL

INTERLEAVING CONCURRENCY

```
interleave: \forall \mu . \mathsf{List}[\{\mathsf{Unit}\}^{\langle \mathsf{async}, \mu \rangle}] \to^{\langle \mathsf{async}, \mu \rangle} \mathsf{Unit}
e_1 \parallel \cdots \parallel e_n \quad \rightsquigarrow \quad \{\mathsf{interleave} \ [e_1, \ldots, e_n]\}
```

PROCESSING MODEL

INTERLEAVING CONCURRENCY

```
interleave: \forall \mu. List[{Unit}^{\langle \mathbf{async}, \mu \rangle}] \rightarrow ^{\langle \mathbf{async}, \mu \rangle} Unit e_1 \parallel \cdots \parallel e_n \quad \rightsquigarrow \quad \{\text{interleave } [e_1, \ldots, e_n]\}
```

The calling context observes effects of the strands!

```
{await x} || {println "foo"} || {yield 1}: {Unit} \( \frac{\text{yield,println,async}}{\} \)
```

```
let rout =
  correlate {
    with h
    x1 from r1
    ...
    xn from rn
    where p
    yield e }
```

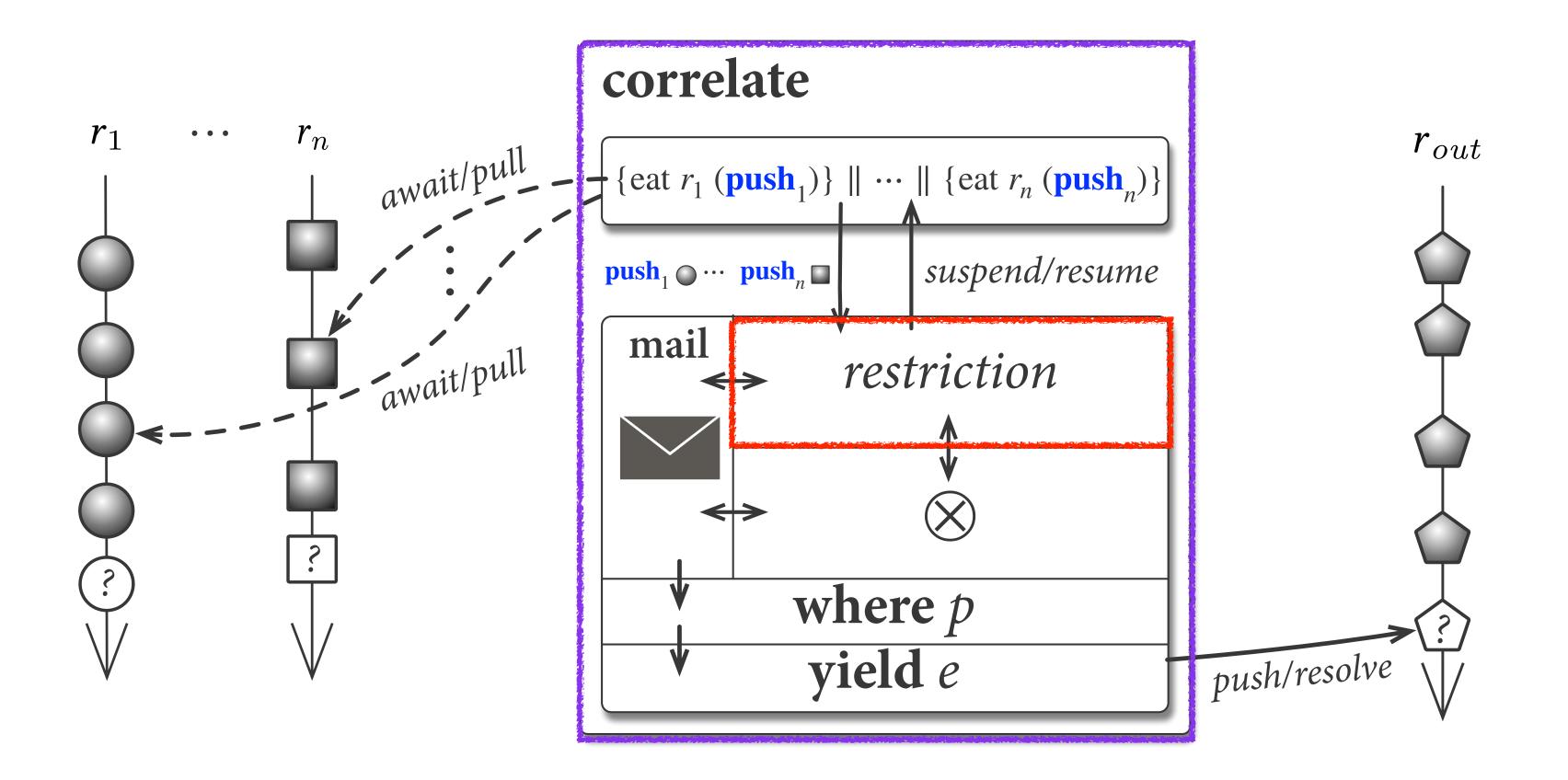
```
let rout =
  correlate {
    with h
    x1 from r1
    ...
    xn from rn
    where p
    yield e }
```

```
let r_{\text{out}} = \text{with } (h_{\otimes} \boxplus h)

\{\text{eat } r_1 \text{ (push}_1)\} \parallel \cdots \parallel \{\text{eat } r_n \text{ (push}_n)\}
```

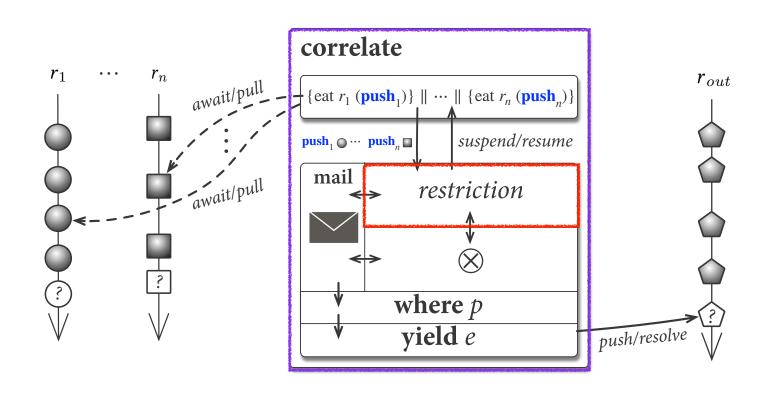
```
let r_{\text{out}} =
correlate {
  with h
  x<sub>1</sub> from r<sub>1</sub>
  ...
  x<sub>n</sub> from r<sub>n</sub>
  where p
  yield e }

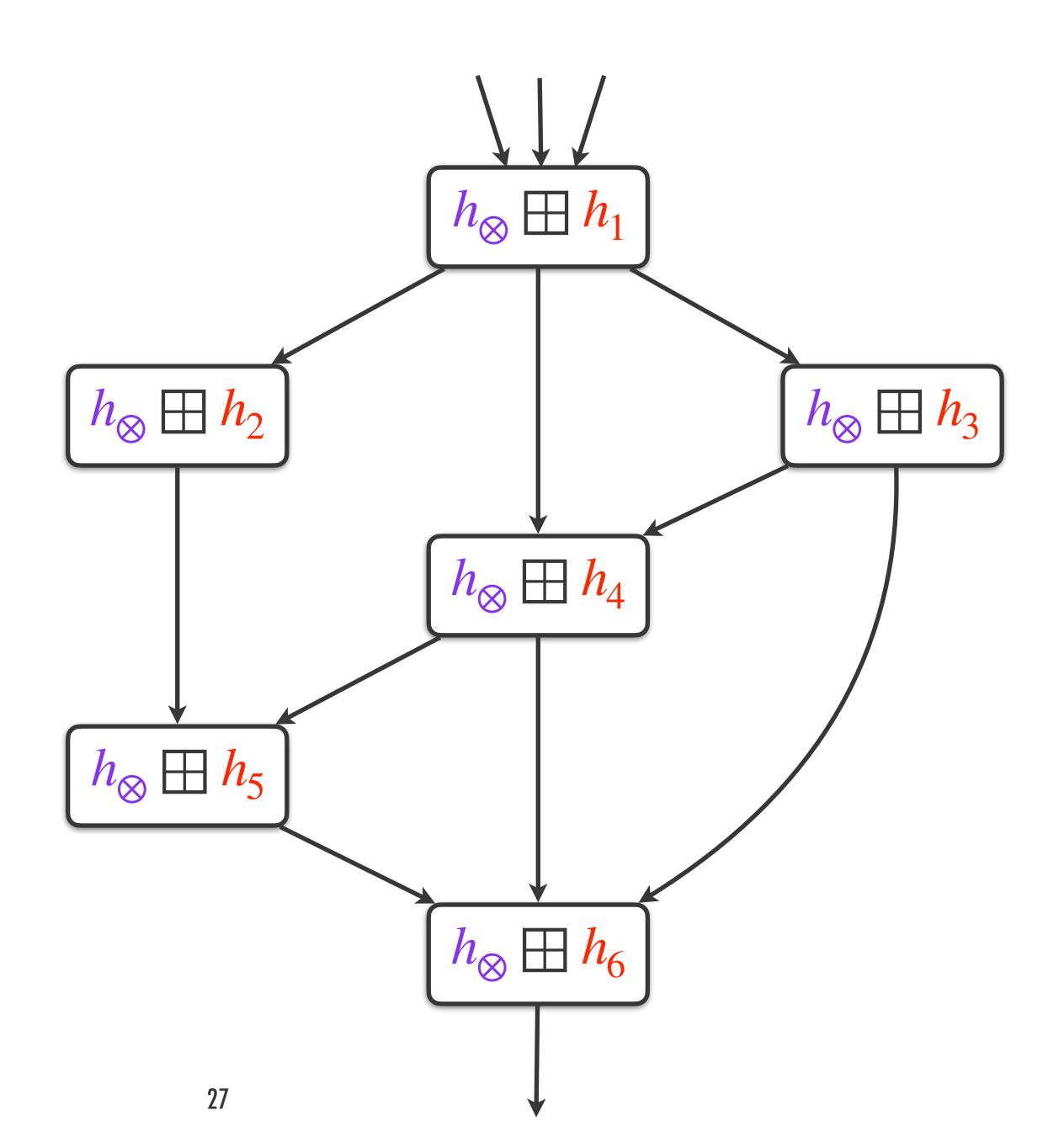
let r_{\text{out}} = \text{with } (h_{\otimes} \boxplus h)
  {eat r_1 (push<sub>1</sub>)} \parallel \cdots \parallel {eat r_n (push<sub>n</sub>)}
```



```
let r_{\text{out}} =
   correlate {
    with h
        x_1 from r_1
        ...
        x_n from r_n
        where p
        yield e }

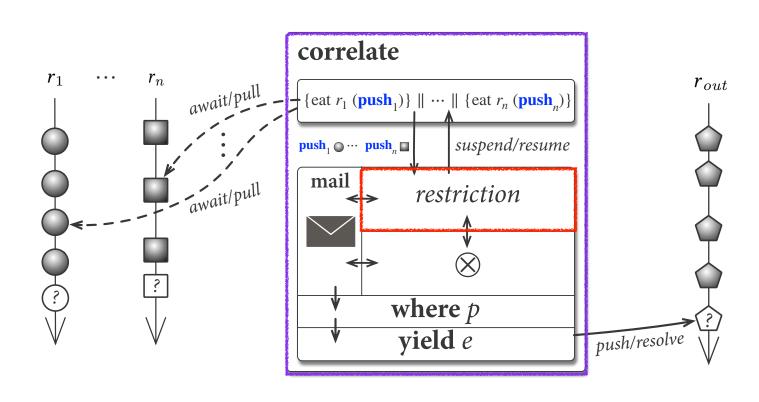
let r_{\text{out}} = \text{with } (h_{\otimes} \boxplus h)
        {eat r_1 (push<sub>1</sub>)} || ... || {eat r_n (push<sub>n</sub>)}
```

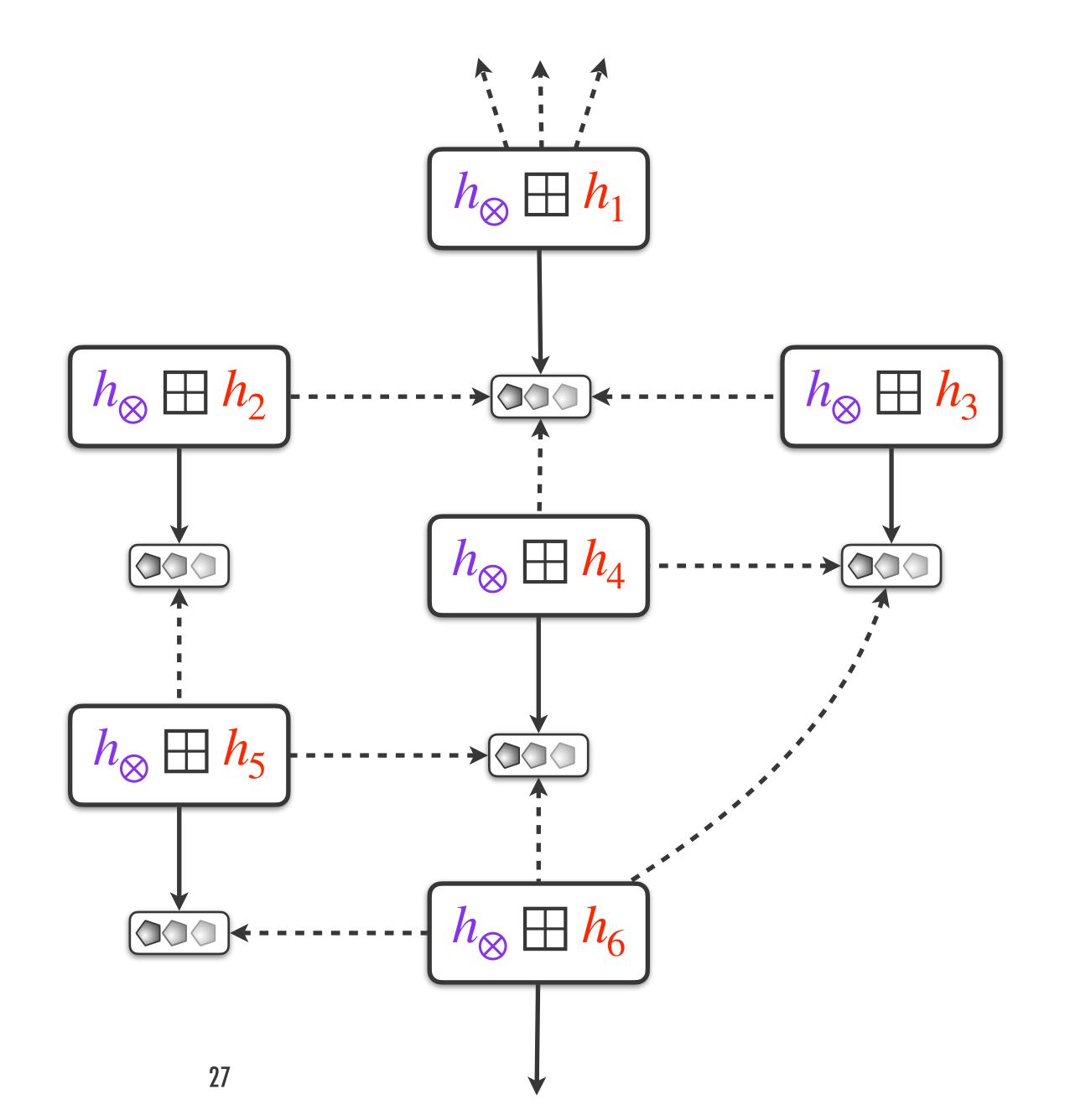




```
let r_{\text{out}} =
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    with h
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    x_n from r_n
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    {eat r_1 (push<sub>1</sub>)} || ... || {eat r_n (push<sub>n</sub>)}
```





NESTING IS BAD!

Could define joins as

$$M[A_1] \rightarrow \ldots \rightarrow M[A_n] \rightarrow M[A_1 \times \ldots \times A_n]$$

return : $A \rightarrow M[A]$

bind:
$$M[A] \rightarrow (A \rightarrow M[B]) \rightarrow M[B]$$

the basis for comprehensions [Wadler '92], LINQ [Meijer '06]:

- bind m_1 ($\lambda x_1:A_1$.
- $\begin{array}{ccc}
 & \text{bind } m_2 \ (\lambda x_2 : A_2.
 \end{array}$

from x_n in m_n

from x_1 in m_1

select
$$\langle x_1, \ldots, x_n \rangle$$

- Problem: sequential binding is inadequate for asynchrony
 - We are not in control of event sources
 - Join should be able to proceed on any event, need non-sequential binding

bind $m_n (\lambda x_n : A_n . \text{return } \langle x_1, \ldots, x_n \rangle) \ldots)$

POLYVARIADICITY

ARITY GENERICITY + HETEROGENEITY, INDEXING BY VARIABLE CONTEXT SHAPE

Pattern Variables in Higher-Order Abstract Syntax (HOAS)

Encodes an indexed function family

$$\mathbf{correlate} \colon \left(\overrightarrow{\mathsf{Shape}[A]} \to (\overrightarrow{A \times \mathsf{Meta}} \to C \times \mathsf{Meta}) \to \mathsf{Shape}[C]\right)_{\overrightarrow{A} \in \mathsf{Type}^*}$$

Payoff: Permits variable binding semantics suitable for data flow patterns, which LINQ and monad comprehensions cannot express!

POLYJOIN: A BETTER LINQ

SUMMARY

TAGLESS INTERPRETERS [CARETTE ET AL. 2009]

- Modular, extensible abstract syntax & interface!
- Polyvariadic binding forms?

POLYVARIADIC FUNCTIONS/PATTERN COMBINATORS [DANVY 1998, RHIGER 2009]

- Separate interface and implementation?
- Polyvariadic binding forms!

POLYJOIN [BRACEVAC ET AL. 2020]

- Modular, extensible abstract syntax & interface!
- Polyvariadic binding forms!

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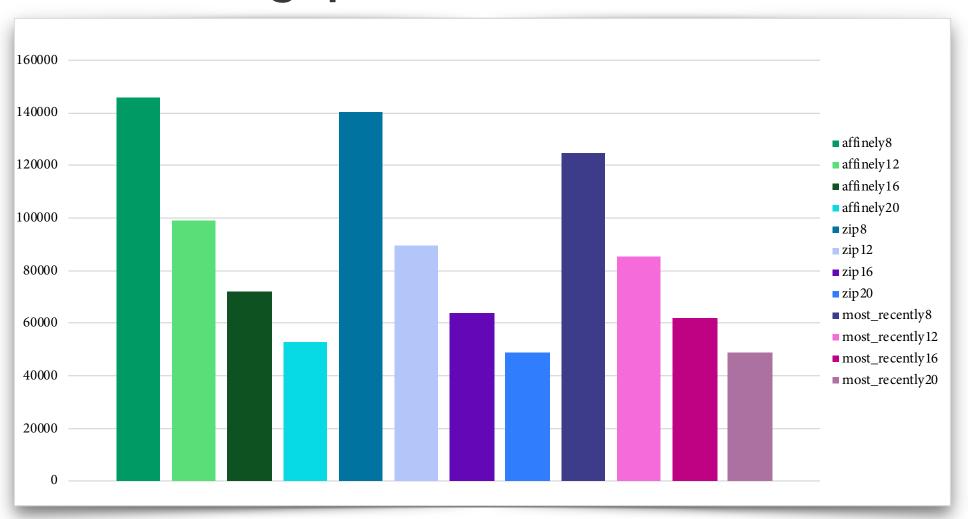
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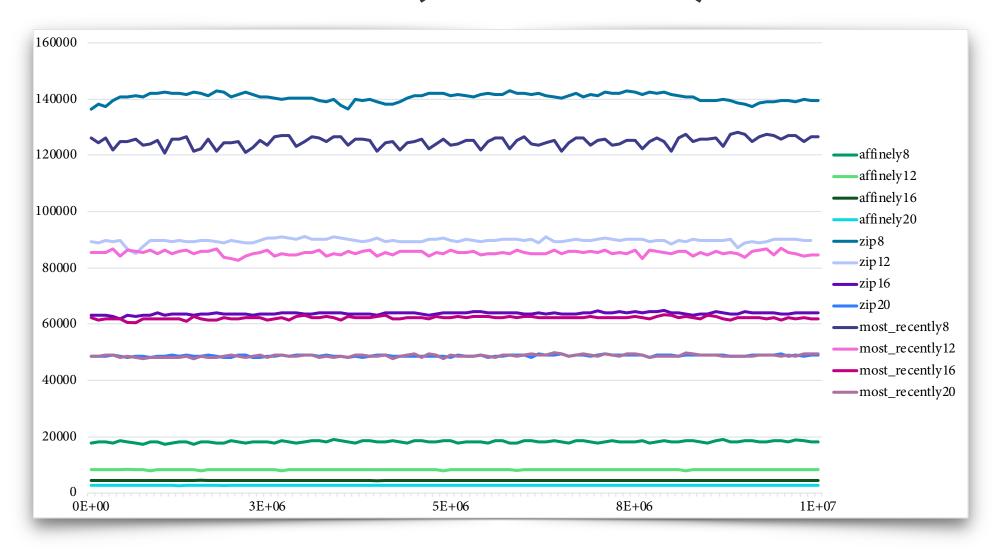
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EFFECT OF ARITY ON PERFORMANCE

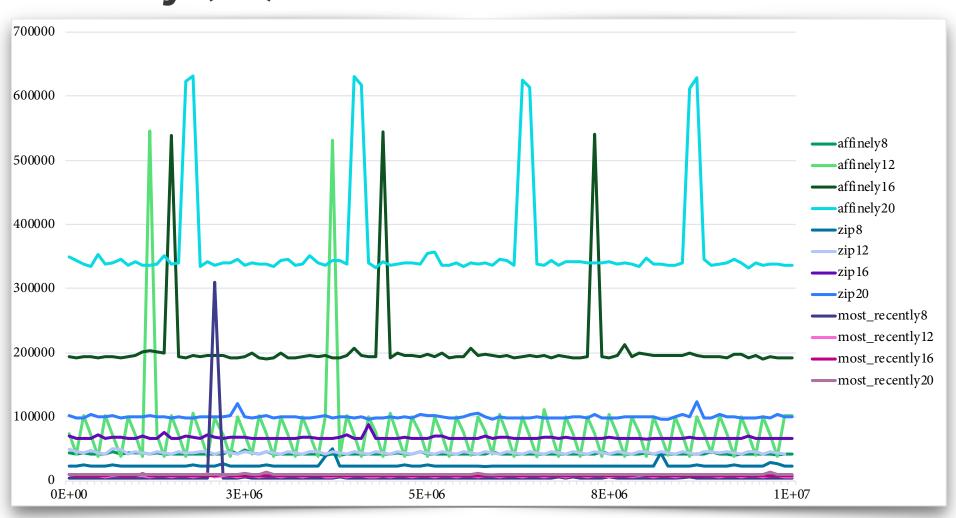
Total Throughput (#Events/sec)



Production Rate (#Events/sec)



Latency (ns)



- Experiment: Uniformly distribute 10^7 events over n inputs, where n = 8,12,16,20.
- Throughput declines with n.
- Culprit: Effects Dispatch/Stacking of handlers.